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Summary: Distribution of a shared symmetric key by a trusted server.
Symmetric key cryptography with server.

Protocol specification (in common syntax)

A, B, S : principal

M, Na, Nb : nonce

Kas, Kbs, Kab : key

1. A → B : M, A, B, {Na, M, A, B}Kas
2. B → S : M, A, B, {Na, M, A, B}Kas , {Nb, M, A, B}Kbs
3. S → B : M, {Na, Kab}Kas, {Nb, Kab}Kbs
4. B → A : M, {Na, Kab}Kas

Description of the protocol rules

The nonce M identifies the session number.

Kas and Kbs are symmetric keys whose values are initially known only by A and S, respectively B and S.

Kab is a fresh symmetric key generated by S in message 3 and distributed to B, directly in message 3, and to A, indirectly, when B forwards blindly {Na, Kab}Kas to A in message 4.

Requirements

The protocol must guaranty the secrecy of Kab: in every session, the value of Kab must be known only by the participants playing the roles of A, B and S.

When A, resp. B, receives the key Kab in message 3, resp. 2, this key must have been issued in the same session by the server S with whom B has started to communicate in message 2.

References

[OR87]

Claimed attacks

Type flaw in [CJ97], where A will accept in last message 4 the triple (M, A, B) as a fresh key Kab.

1. A -> I(B) : M, A, B, {Na, M, A, B}Kas
2. B -> S : M, A, B, {Na, M, A, B}Kas , {Nb, M, A, B}Kbs
3. S -> B : M, {Na, Kab}Kas, {Nb, Kab}Kbs
4. I(B) -> A : M, {Na, M, A, B}Kas

Citations

[CJ97] John Clark and Jeremy Jacob. A survey of authentication protocol literature : Version 1.0., November 1997.

[OR87] D. Otway and O. Rees. Efficient and timely mutual authentication. *Operating Systems Review*, 21(1):8–10, 1987.